## Page Tables

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Lecture 12
2.0.0

## What Should We Do About the Page Table Size?

Most programs don't use all the virtual memory space, how can we take advantage?

## We Can Make Our Page Table Fit on a Page



## Multi-Level Page Tables Save Space for Sparse Allocations



## Page Allocation Uses A Free List

Given physical pages, the operating system maintains a free list (linked list)
The unused pages themselves contain the next pointer in the free list Physical memory gets initialized at boot

To allocate a page, you remove it from the free list
To deallocate a page you add it back to the free list

## Insight: Use a Page for Each Smaller Page Table

There are $512\left(2^{9}\right)$ entries of 8 bytes $\left(2^{3}\right)$ each, which is 4096 bytes
The PTE for $\mathrm{L}(\mathrm{N})$ points to the page table for $\mathrm{L}(\mathrm{N}-1)$
You follow these page tables until LO and that contains the PPN

## The Smaller Page Tables are Just Like Arrays

Instead of:
int page_table[512] // What's the size of this?
or

```
x = page_table[2]; // What's the offset of index 2?
```

You have:
PTE page_table[512]
where:
sizeof(page_table) $==$ PAGE_SIZE
and
sizeof(page_table) $=$ number of entries * sizeof(PTE)

## Consider Just One Additional Level

Assume our process uses just one virtual address at 0x3FFFF008 or 0b11_1111_1111_1111_1111_0000_0000_1000 or Ob111111111_111111111_000000001000

We'll just consider a 30-bit virtual address with a page size of 4096 bytes We would need a 2 MiB page table if we only had one $\left(2^{18} \times 2^{3}\right)$

Instead, we have a 4 KiB L1 page table $\left(2^{9} \times 2^{3}\right)$ and a 4 KiB L0 page table Total of 8 KiB instead of 2 MiB

Note: worst case if we used all virtual addresses we would consume $2 \mathrm{MiB}+$ 4 KiB

## Translating 3FFFF008 with 2 Page Tables

Consider the L1 table with the entry:

| Index | PPN |
| ---: | ---: |
| 511 | $0 \times 8$ |

Consider the L0 table located at $0 \times 8000$ with the entry:
Index PPN

511 DxCAFE
The final translated physical address would be: 0xCAFEOO8

